

Chapter 04: Relational Algebra

1. Introduction

The first language covered in this document is **relational algebra**. It consists of a set of operations used to manipulate relations, which are considered as sets of tuples.

In relational algebra, all operations manipulate one or two relations, and the final result is also a relation.

Generally, there are two types of operations:

↳ **Set operations:**

- Union
- Intersection
- Difference

↳ **Specific operations:**

- Projection
- Selection
- Join
- Division

Set Operations

2.1 Union

The union of two relations R_1 and R_2 with the same schema is a relation R_3 with an identical schema that contains the tuples present in R_1 , R_2 , or both.

Conditions for Union:

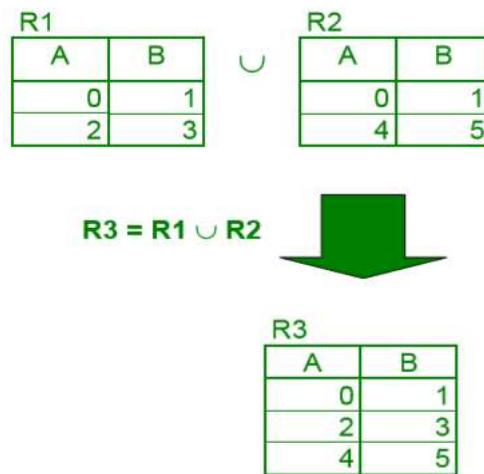
- ✓ **Same schema:**

- The attributes (columns) in **R1** and **R2** must have the same number, order, and domain (data type).

✓ **No duplicates:**

- If a tuple exists in both **R1** and **R2**, it appears only once in the result.

Example 1:



Syntax : Let $R1, R2$ two relation with the same schema, the union is noted:

$$R1 \cup R2$$

Example 2:

Relations:

R1 (Students in Class A):

Name	Age
Alice	20
Bob	22

R2 (Students in Class B):

Name	Age
Charlie	21

Name	Age
Alice	20

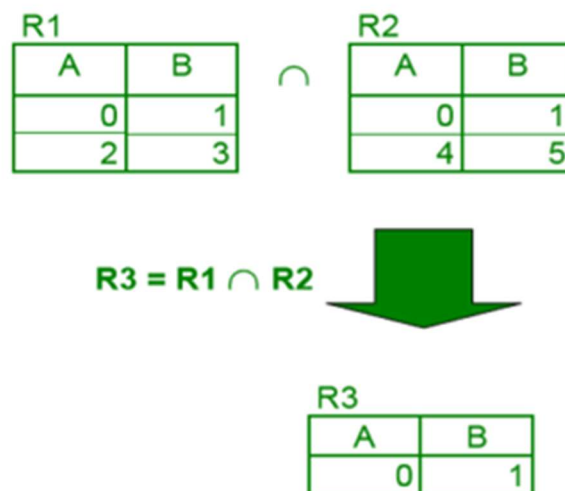
Union
 $R1 \cup R2$

Result:

Name	Age
Alice	20
Bob	22
Charlie	21

2.2 Intersection:

The intersection of two relations $R1$ and $R2$ with the same schema is a relation $R3$ with an identical schema, containing the tuples that are common to both $R1$ and $R2$.



Syntax : Let $R1, R2$ two relation with the same schema, the intersection is noted:

$$R1 \cap R2$$

Example 04:

Relations:

$R1$ (Students enrolled in Math):

Name	Age
Alice	20
Bob	22

R2 (Students enrolled in Physics):

Name	Age
Alice	20
Charlie	21

Intersection
 $R1 \cap R2$

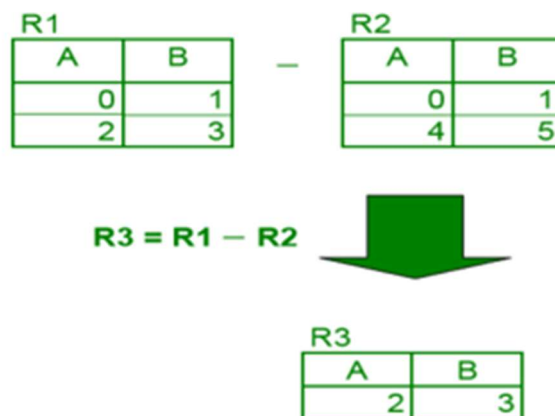
Result:

Name	Age
Alice	20

2.3 Difference:

The difference between two relations $R1$ and $R2$ with the same schema is a relation $R3$ with an identical schema, containing the tuples from $R1$ that do not belong to $R2$.

Example 05:



Syntax: Let $R1, R2$ two relation with the same schema, the difference is noted:

$$R1 - R2$$

Example 06

Relations:

R1 (Students enrolled in Math):

<i>Name</i>	<i>Age</i>
Alice	20
Bob	22
Charlie	21

R2 (Students enrolled in Physics):

<i>Name</i>	<i>Age</i>
Alice	20
Charlie	21

Difference operation

$R1 - R2$

Result:

The result is a relation $R3$, containing the tuples from $R1$ that are **not** in $R2$.

R3 (Students in Math but not in Physics):

<i>Name</i>	<i>Age</i>
Bob	22

✓ **Specific Operations:**

3.1 Projection:

⚡ The projection of a relation $R1$ is the relation $R2$ obtained by removing the attributes of $R1$ that are not mentioned and then eliminating any duplicate tuples if necessary.

syntax : $\pi A1, A2, \dots, An (R1)$

⚡ The projection of a relation $R1$ is performed on the attributes $A1, A2, \dots, An$.

⚡ Projection corresponds to a **vertical slicing**.

Example 07:

$\pi_{A1, A3}(R)$

A1	A2	A3	A4

Example 08:

Let Product Table Example:

N°Product	Name	Price	QStock
01	car hyundai eon	6000 euro	10
02	Motocycle MBK	10000 euro	20
03	bicycle Detwin	300 euro	30

For each product give the quantity in stock:

$\pi_{Name, Qstock}(Product)$

Result:

Name	QStock
car hyundai eon	10
Motocycle MBK	20
bicycle Detwin	30

3.2 Restriction:

⚡ The restriction of a relation $R1$ is a relation $R2$ with the same schema, containing only the tuples from $R1$ that satisfy the specified condition.

syntax: $\sigma_{condition}(R1)$.

⚡ The restriction of a relation $R1$ based on the "condition" criterion, where "condition" is a comparison relation between two attributes or between an attribute and a value.

⚡ The restriction operation allows the extraction of tuples that satisfy a condition.

⚡ The restriction corresponds to a **horizontal slicing** of the table that describes the relation.

Example 09:

	A1	A2	A3	A4
→				
→				

Example 10: give all product where the quantity in stock superior 10.

$\sigma_{Qstock > 10}(\text{Product})$.

3.3 Join:

✓ The join of two relations $R1$ and $R2$ is a relation $R3$, where the tuples are obtained by concatenating the tuples from $R1$ with those from $R2$, keeping only those that satisfy the join condition.

syntax: $R1 \bowtie R2(\text{condition})$.

✓ **Steps of join**

1. **Cartesian Product:** First, perform the Cartesian product between **R1** and **R2**. This generates all possible combinations of tuples between **R1** and **R2**.

2. **Restriction:** Then, apply the join condition to keep only the tuples that satisfy the condition.

- ✓ The tuples of $R1 \bowtie R2(\text{condition})$ are all pairs $(u1, u2)$, where $u1$ is a tuple from $R1$ and $u2$ is a tuple from $R2$ that satisfy the condition.
- ✓ The join of two relations $R1$ and $R2$ is the **Cartesian product** of the two relations followed by a restriction based on the condition.

The join condition must be of the form: $::$

Where: $\text{attribute1} \in 1\text{st relation}$ and $\text{attribute2} \in 2\text{nd relation}$, and $"::"$ is a comparison operator.

Remark:

- The join allows combining two relations using a join condition.
- Normalization (putting into canonical normal form) leads to decomposing the relation.

Example 11:

Let the following two relation:

Product(Prod-ID, Name, Uprice)

Sale (ID-client, Prod-ID, date, QTE-sold)

Give for each sale the product reference, its description, price, ID client, date, and quantity sold.

The table product:

Prod-ID	Name	UPrice
01	car hyundai eon	6000 euro
02	Motocycle MBK	10000 euro
03	bicycle Detwin	300 euro

The table sale:

ID-client	Prod-ID	Date	QTE-sold
01	01	01/01/1998	1
02	02	02/01/1998	3
03	01	03/01/1998	1

Product \bowtie Sale (Product.Prod-ID=Sale.Prod-ID)

Prod-ID	Name	UPrice	ID-client	Date	QTE-sold
01	car hyundai eon	6000 euro	01	01/01/1998	1
02	Motocycle MBK	10000 euro	02	02/01/1998	3
01	car hyundai eon	6000 euro	03	03/01/1998	1

3.4 Division:

Let there be two relations $R1(A1, A2, \dots, An, B1, B2, \dots, Bm)$, $R2(B1, B2, \dots, Bm)$ if the schema of $R2$ is a sub-schema of $R1$, the division of $R1$ by $R2$ is a relation $R3$ where:

- The schema of $R3$ is the complementary sub-schema of $R2$ with respect to $R1$.
- A tuple $(a1, a2, \dots, an)$ belongs to $R3$ if $(a1, a2, \dots, an, b1, b2, \dots, bm)$ belongs to $R1$ for all $(b1, b2, \dots, bm) \in R2$.

Syntax: $R1 \div R2$

Example 12:

Here are the following tables; which students are enrolled in all sports?

Here the table Enrollment

Student	sport
julie	Danse
jean	judo
julie	Hand-ball
julie	judo

here the table Sport

sport
Danse
judo
Hand-ball

Student ÷ Sport

Student
julie

3.5 Aggregation:

In Relational Algebra, γ (gamma) represents the **aggregation (grouping) operator**. It is used to perform **grouping and aggregate functions** like:

- **SUM** (Total)
- **AVG** (Average)
- **COUNT** (Number of occurrences)
- **MIN** (Minimum)
- **MAX** (Maximum)

Syntax:

γ grouping attributes, aggregate function(attribute)(relation)

Example 13:

New Database Schema: E-Commerce System

customers (customer_id, name, email, city)

orders (order_id, customer_id, order_date, total_amount)

products (product_id, product_name, price)

order_Items (order_item_id, order_id, product_id, quantity)

Example Queries using γ (Gamma) in Relational Algebra:

1. Find the total number of orders per customer

γ customer_id, COUNT(order_id) (orders)

Explanation:

- Groups the orders table by customer_id.
- Counts how many orders each customer has placed.

2. Find the total revenue (sum of all orders) per customer

γ customer_id, SUM(total_amount)(orders)

Explanation:

- Groups the orders table by customer_id.
- Computes the **total spending** of each customer.

3. Find the most expensive product in each order

γ order_id, MAX(price) (Order_items \bowtie products (order_items.product_id=products.product_id))

Explanation:

- Joins order_items and products to get the price of each product in an order.
- Groups by order_id and finds the **highest price** product in each order.

4. Find the average quantity of items per order

γ order_id, AVG(quantity)(order_items)

Explanation:

- Groups the order_items table by order_id.
- Computes the **average quantity** of items purchased per order.

5. Find the total quantity of each product sold

γ product_id, SUM(quantity)(order_items)

Explanation:

- Groups order_items by product_id.

- Computes the **total quantity** sold for each product.

4. Some examples in Algebraic Language:

The algebraic language allows formulating a question through a series of relational algebra operations. This question is called a query.

Example 14:

Queries on the CLIENT, PRODUCT, SALE schema of Computer hardware sales office

CLIENT (IdCli, name, city)

PRODUCT (IdPro, description, brand, price)

SALE (IdCli, IdPro, date, quantity)

Question :

1. Give the client live in constantine

$\sigma (\text{city} = \text{'constantine'}) (\text{CLIENT})$

2. Give the product numbers of Apple brand and with a price < 50000 DA.

$\pi \text{IdPro} (\sigma (\text{brand} = \text{'Apple'} \wedge \text{price} < 50000) (\text{PRODUCT}))$

3. Give the brand and the price of the product sold 01/01/2020 with the quantity > 3

$\pi \text{brand, price} (\sigma (\text{date} = \text{'01/01/2020'} \text{ AND } \text{quantity} > 3) (\text{PRODUCT} \bowtie \text{SALE} (\text{PRODUCT.IdPro} = \text{SALE.IdPro})))$

Example 15:

Let the following schema of database of the application Diabetes.AI

user (User-ID, Lastname-user, Firstname-user, email, pass-word, DOB, diabetes_type, Date-creation)

blood-Sugar-measurement (Measurement-ID, User-ID, value, measurement-date, time-of-day)

meal_logs (Meal-ID, User-ID, photo-url, description, glycemic-index, recorded-at)

chatbot-interactions (Intercation-ID, User-ID, question, reponse, date-interaction)

Questions:

Q1) Retrieve all blood sugar readings for a specific user ($user_id = 1$).

$\sigma_{User-ID=1}(blood_sugar_measurement)$

Q2) Find meals with a high glycemic index (>70).

$\sigma_{glycemic_index>70}(meal_logs)$

Q3) Get all chatbot interactions for a specific user ($user_id = 1$).

$\sigma_{User-ID=1}(chatbot_interactions)$

Q4) Retrieve all users who have Type 2 diabetes.

$\sigma_{diabetes_type='Type2'}(users)$

Q5) Get the names of users along with their latest blood sugar measurement.

$\pi_{first_name, last_name, value, measurement_date}(users \bowtie_{User-ID=blood_sugar_measurement.User-ID} blood_sugar_measurement)$

Q6) List meals eaten by a user along with their blood sugar levels after eating.

$\pi_{meal_logs.description, blood_sugar_measurement.value,$

$blood_sugar_measurement.measurement_date}((meal_logs \bowtie_{meal_logs.User-ID=blood_sugar_measurement.User-ID} blood_sugar_measurement$

$(meal_logs.User-ID=blood_sugar_measurement.User-ID))$

Q7) Retrieve all meals recorded by a specific user ($user_id = 2$).

$\sigma_{User-ID=2}(meal_logs)$

Q8) Find users who have recorded at least one blood sugar reading.

$\pi_{users.User-ID, users.first_name, users.last_name}(users \bowtie_{users.User-ID=blood_sugar_measurement.User-ID} blood_sugar_measurement)$

Q9) Retrieve all chatbot interactions where the response contains the word "insulin".

$\sigma_{response LIKE \% 'insulin' \%}(chatbot_interactions)$

Q10) Find the average blood sugar level of a specific user (User_ID = 3).

γ user_id, AVG(value)(σ User-ID=3(blood_sugar_measurement))

Q11) Get the max blood sugar reading for specific user (User_ID = 3).

γ user_id, MAX(value)(σ User-ID=3(blood_sugar_measurement))

Q12) Retrieve all users who have never interacted with the chatbot.

π users.User-ID, users.first_name, users.last_name(users) -

π users.User-ID, users.first_name, users.last_name(users \bowtie chatbot_interactions(users.User-ID = chatbot_interactions.User-ID))

Q13) Retrieve meals with a high glycemic index (>70) eaten by users with Type 2 diabetes.

π meal_logs.description, meal_logs.glycemic_index, users.first_name, users.last_name(σ meal_logs.glycemic_index > 70 (meal_logs \bowtie users)(meal_logs.User-ID = users.User-ID \wedge users.diabetes_type = 'Type2'))

Q14) Find the total number of blood sugar readings recorded per user.

γ User-ID, COUNT(User-ID)(blood_sugar_measurement)

Q15) Retrieve the users who have logged both a meal and a blood sugar reading on the same day.

π users.User-ID, users.first_name, users.last_name((users \bowtie meal_logs \bowtie blood_sugar_measurement)((meal_logs.User-ID = users.User-ID) \wedge (meal_logs.User-ID = blood_sugar_measurement.User-ID \wedge meal_logs.recorded-at = blood_sugar_measurement.measurement-date)))

Q16) Find the highest glycemic index recorded per user.

γ User-ID, MAX(glycemic_index)(meal_logs)

Q17) Find the total number of chatbot interactions per user

γ User-ID, COUNT(Intercation-ID) chatbot_interactions