

Solution Tutorial 02 :

Exercise 01 :

We have in the following functional dependencies C dependent only from A:

$A \rightarrow E, F, G, C$;

$A, B \rightarrow C, M$; $\neg 2NF$

$B \rightarrow I, J$;

$E \rightarrow G, H$;

$J \rightarrow K$

So after the normalization to 2NF we have :

$A \rightarrow E, F, G, C$; $\neg 3NF$

$A, B \rightarrow M$;

$B \rightarrow I, J$;

$E \rightarrow G, H$;

$J \rightarrow K$

Now the normalization to 3NF: we have G dependent to E so we eliminate the transitivity.

$A \rightarrow E, F, C$;

$A, B \rightarrow M$;

$B \rightarrow I, J$;

$E \rightarrow G, H$;

$J \rightarrow K$

Now all the FD are in 3NF.

Exercise 02: Learn how to deduce transitivity between FDs:

Q1. $\{x \rightarrow y ; z \rightarrow w\} \Rightarrow xz \rightarrow yw$ oui $x \times z \rightarrow y \times w$ by concatenation.

Q2. $\{xy \rightarrow z ; z \rightarrow x\} \Rightarrow z \rightarrow y$ non

Lastname, Firstname \rightarrow address

address \rightarrow Lastname but address $\not\rightarrow$ Firstname in a house can live many persons.

Q3. $\{x \rightarrow y ; y \rightarrow z\} \Rightarrow x \rightarrow yz$ yes by transitivity on a $x \rightarrow y$ et $y \rightarrow z$ so $x \rightarrow z$ which means $x \rightarrow y$ and z hence we deduce that $x \rightarrow yz$

Q4. $\{x \rightarrow y ; w \rightarrow z\}$ et $w \supseteq y \Rightarrow x \rightarrow z$ non

Counterexample :

Address \rightarrow Lastname

Lastname, Firstname \rightarrow age but address $\not\rightarrow$ age.

Q5. $\{w \rightarrow y, x \rightarrow z\} \Rightarrow wx \rightarrow y$ yes by concatenation $wx \rightarrow y$ and z ce qui signifie que $wx \rightarrow y$ and $wx \rightarrow z$

Q6. $\{x \rightarrow y\}$ et $y \supseteq z \Rightarrow x \rightarrow z$ yes $x \rightarrow y$ and $y \supseteq z$ so $y = z \cup w$ such that $w = y - z$ which means $x \rightarrow y$ we have $x \rightarrow z \cup w$ so $x \rightarrow z$ and $x \rightarrow w$.

Which means $x \rightarrow y$ so $x \rightarrow z \cup w$ hence $x \rightarrow z$ and $x \rightarrow w$

Q7. $\{x \rightarrow y, x \rightarrow w, wy \rightarrow z\} \Rightarrow x \rightarrow z$ yes

$x \rightarrow y$ and $x \rightarrow w$ and we have $wy \rightarrow z$ which means $x \rightarrow z$.

Q8. $\{x, y \rightarrow z, y \rightarrow w\} \Rightarrow xw \rightarrow z$ no counterexample :

Lastname, N°mobil \rightarrow Firstname

N°mobil \rightarrow address mais nom, address \rightarrow Firstname

Q9. $\{x \rightarrow y, xy \rightarrow z\} \Rightarrow x \rightarrow z$ yes $x \rightarrow y$ and $xy \rightarrow z$ by transitivity we obtain y from x so we can eliminate from $xy \rightarrow z$ we obtain $x \rightarrow z$.

Exercise 03 :

1. Universal schema:

$R(\text{Book-ID}, \text{Title}, \text{Lname-A}, \text{Fname-A}, \text{Copy-ID}, \text{Nbre-pages}, \text{Code-Pub}, \text{Name-Pub}, \text{year}, \text{Code-depot}, \text{Name-depot}, \text{City}, \text{Date}, \text{Qte})$

2. Normalization:

$R1(\text{Book-ID}, \text{Title}, \text{Lname-A}, \text{Fname-A}) \rightarrow 2NF$

$R2(\text{Book-ID}, \text{Copy-ID}, \text{Nbre-pages}, \text{Code-Pub}, \text{Name-Pub}, \text{year}) \rightarrow 2NF$

$R3(\text{Book-ID}, \text{Copy-ID}, \text{Code-depot}, \text{Name-depot}, \text{City}, \text{Date}, \text{Qte}) \rightarrow 2NF$

Normalization to 2NF

$R1(\text{Book-ID}, \text{Title})$

$R'_1(\text{Book-ID}, \text{Lname-A}, \text{Fname-A})$

$R2(\text{Book-ID}, \text{Copy-ID}, \text{Nbre-pages})$

$R'_2(\text{Book-ID}, \text{Copy-ID}, \text{Code-Pub}, \text{year})$

$R''_2(\text{Code-Pub}, \text{Name-Pub})$

$R3(\text{Code-depot}, \text{Name-depot}, \text{City})$

$R'_3(\text{Book-ID}, \text{Copy-ID}, \text{Code-depot}, \text{Date}, \text{Qte})$

All relation now are in 2NF

Normalization to 3NF

All relation now are in 3NF

Renaming relations:

Book (Book-ID, Title)

Author (Book-ID, Lname-A, Fname-A)

Copy(Book-ID, Copy-ID, Nbre-pages)

Published (Book-ID, Copy-ID, Code-Pub, year)

Publisher (Code-Pub, Name-Pub)

Depot(Code-depot, Name-depot, City)

Stored (Book-ID, Copy-ID, Code-depot, Date, Qte)

Exercise 04:

1. Universal schema:

$R(\underline{Supp-ID}, Name - Supp, Adr-Supp, \underline{Prod-ID}, Prod-Des, Prod-name, Weight, Unit-Price, \underline{Lname-Mer}, \underline{Fname-Mer}, Adr-Mer, Date-order, QTE-purshased, Date-delivery, QTE-delivered)$ 1NF but $\neg 2NF$

2. Normalization:

$R1(\underline{Supp-ID}, Name - Supp, Adr-Supp, \underline{Prod-ID}) \neg 2NF$

$R2(\underline{Prod-ID}, Prod-Des, Prod-name, Weight, Unit-Price,)$

$R3(\underline{Lname-Mer}, \underline{Fname-Mer}, Adr-Mer)$

$R4(\underline{Supp-ID}, \underline{Lname-Mer}, \underline{Fname-Mer}, \underline{Prod-ID}, \underline{Date-order}, \underline{QTE-purshased})$

$R5(\underline{Supp-ID}, \underline{Lname-Mer}, \underline{Fname-Mer}, \underline{Prod-ID}, \underline{Date-delivery}, \underline{QTE-delivered})$

Put it in 2NF:

$R'_1(\underline{Supp-ID}, Name - Supp, Adr-Supp)$

$R''_1(\underline{Supp-ID}, \underline{Prod-ID})$

$R2(\underline{Prod-ID}, Prod-Des, Prod-name, Weight, Unit-Price,)$

$R3(\underline{Lname-Mer}, \underline{Fname-Mer}, Adr-Mer)$

$R4(\underline{Supp-ID}, \underline{Lname-Mer}, \underline{Fname-Mer}, \underline{Prod-ID}, \underline{Date-order}, \underline{QTE-purshased})$

$R5(\underline{Supp-ID}, \underline{Lname-Mer}, \underline{Fname-Mer}, \underline{Prod-ID}, \underline{Date-delivery}, \underline{QTE-delivered})$

All relation in 2NF and in 3NF.

Renaming relation:

Supplier(Supp-ID, Name - Supp, Adr-Supp)

Have(Supp-ID, Prod-ID)

Product(Prod-ID, Prod-Des, Prod-name, Weight, Unit-Price,)

Merchant(Lname-Mer, Fname-Mer, Adr-Mer)

Order(Supp-ID, Lname-Mer, Fname-Mer, Prod-ID, Date-order, QTE-purchased)

Delivery(Supp-ID, Lname-Mer, Fname-Mer, Prod-ID, Date-delivery, QTE-delivered)

Exercise 05:

1. Universal schema:

R(Stud-ID, Stud-Lname, Stud-Fname, Stud-Adr, Course-C, Title, Summary, Code-pgm, Name-pgm, Prof-ID, Prof-Lname, Prof-Fname, Rank, Prof-Adr) 1NF but \neg 2NF

2. Normalization:

Step01:

*R*₁(Stud-ID, Stud-Lname, Stud-Fname, Stud-Adr, code-pgm) \neg 2NF

*R*₂(Prof-ID, Code-pgm, Name-pgm) \neg 2NF

*R*₃(Stud-ID, Course-C, title, summary) \neg 2NF

*R*₄(Prof-ID, Prof-Lname, Prof-Fname, Rank, Prof-Adr, Course-C) \neg 2NF

Step02:

R'₁(Stud-ID, Stud-Lname, Stud-Fname, Stud-Adr) 2NF

R''₁(Stud-ID, Code-pgm) 2NF

R'₂(Code-pgm, Name-pgm)

R''₂(Prof-ID, Code-pgm)

R'₃(Course-C, title, summary)

R''₃(Stud-ID, Course-C)

R'₄(Prof-ID, Prof-Lname, Prof-Fname, Rank, Prof-Adr)

R''₄(Prof-ID, Course-C)

Step 03:

All relation in 3NF.

Renaming relation:

Student(Stud-ID, Stud-Lname, Stud-Fname, Stud-Adr) 3NF

Choose(Stud-ID, Code-pgm) 3NF

Program(Code-pgm, Name-pgm) 3NF

Managed(Prof-ID, Code-pgm) 3NF

Course(Course-C, Title, Summary) 3NF

Take(Stud-ID, Course-C) 3NF

Professor(Prof-ID, Prof-Lname, Prof-Fname, Rank, Prof-Adr) 3NF

Created(Course-C, Prof-ID) 3NF